

The Power of Hard Attention Transformers on Data Sequences: A Formal Language Theoretic Perspective

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Motivation

- Transformers are the basic model in machine learning used in recent LLMs



- Sometimes their output is just wrong:

Could you please count the number of "R"s in the word "arbitrary"?

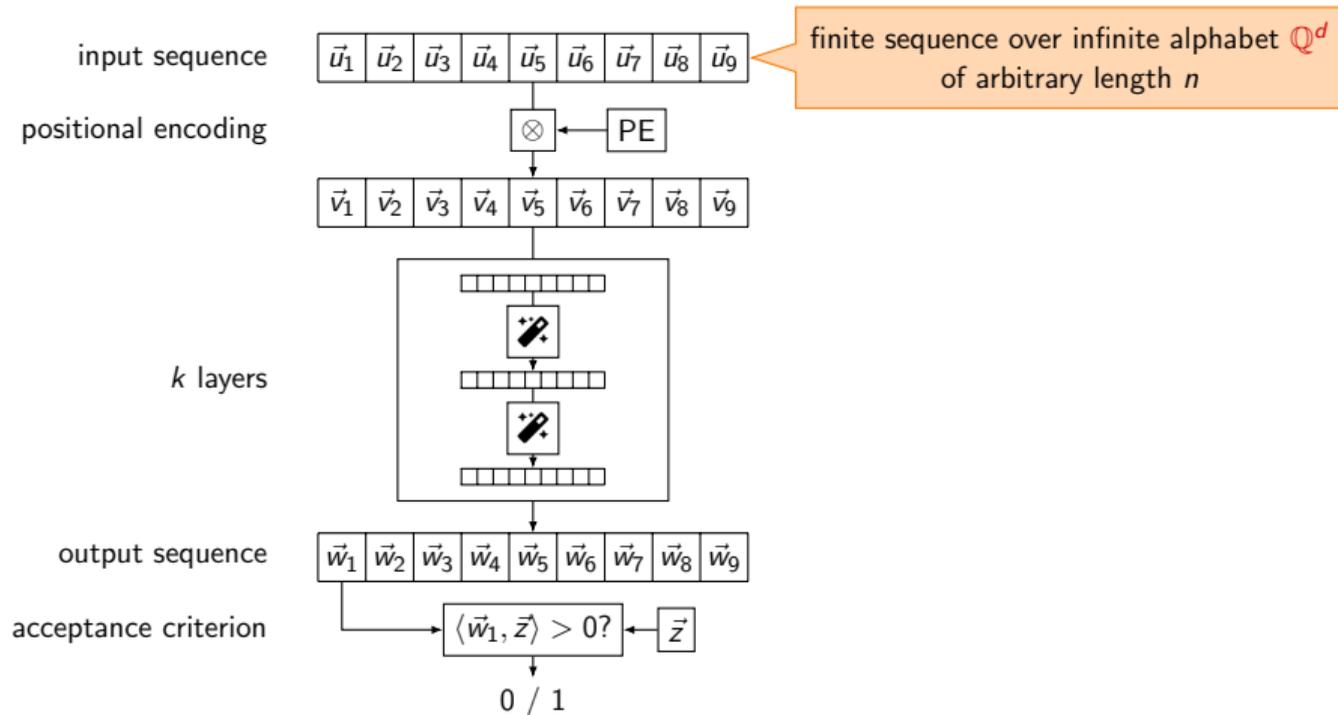


The word "arbitrary" contains **2** "R"s.

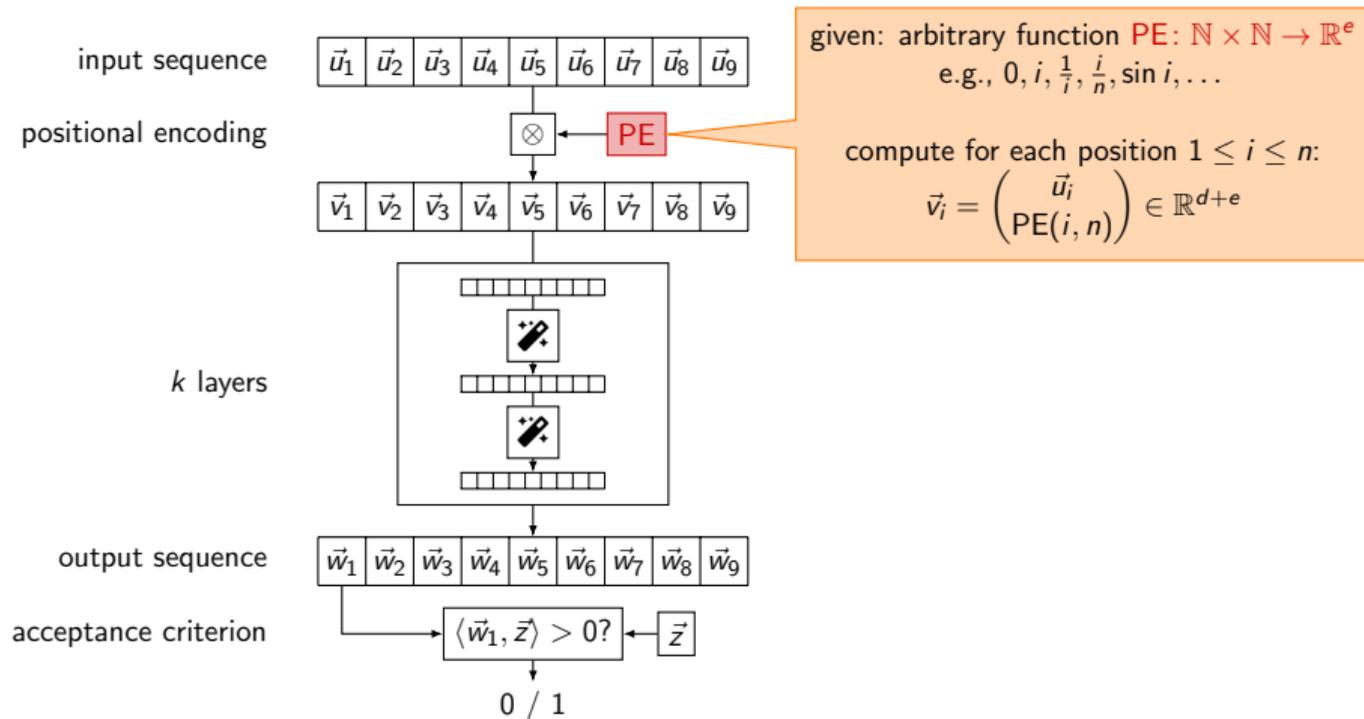
(Spoiler: "arbitrary" contains **three** "R"s)

- Here: we analyze expressibility of transformers that take complex data (like pictures, videos, voice, time series, ...) as input

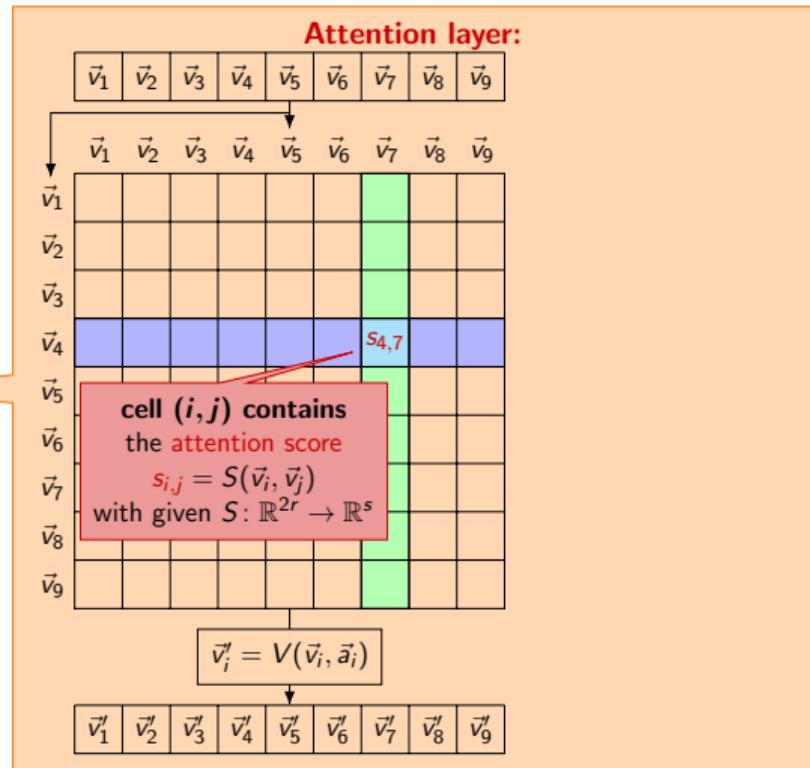
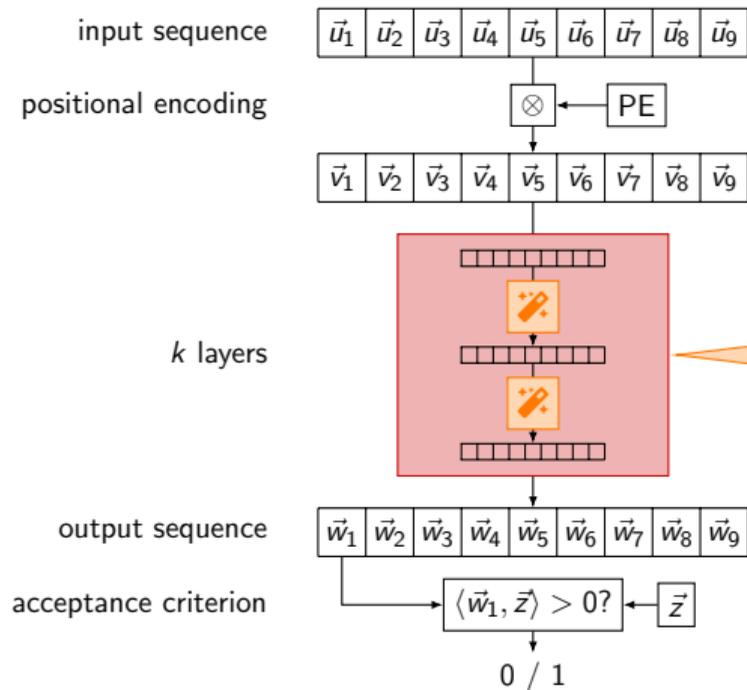
Unique Hard Attention Transformers (UHAT) on Data Sequences



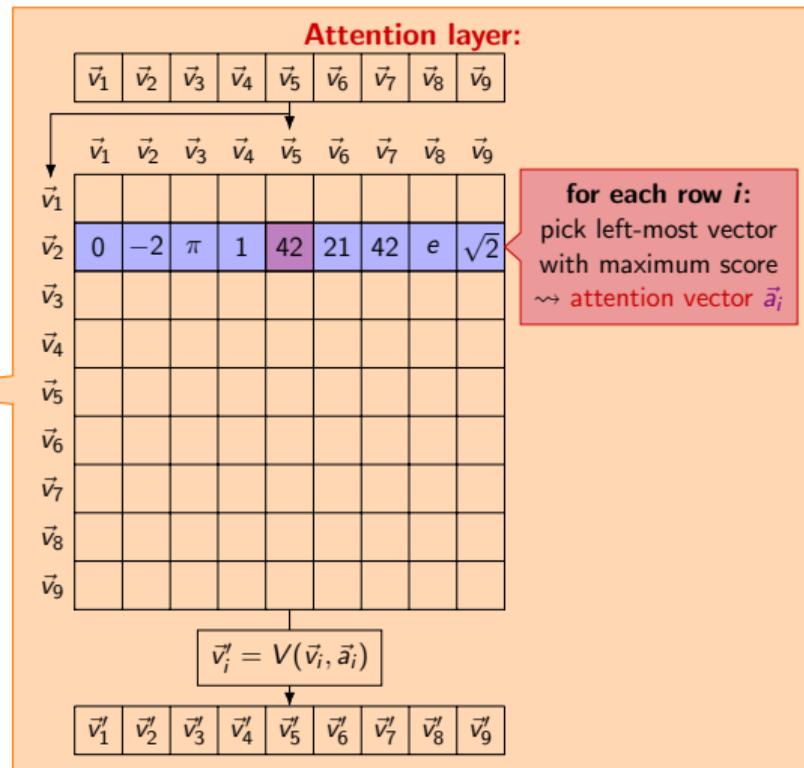
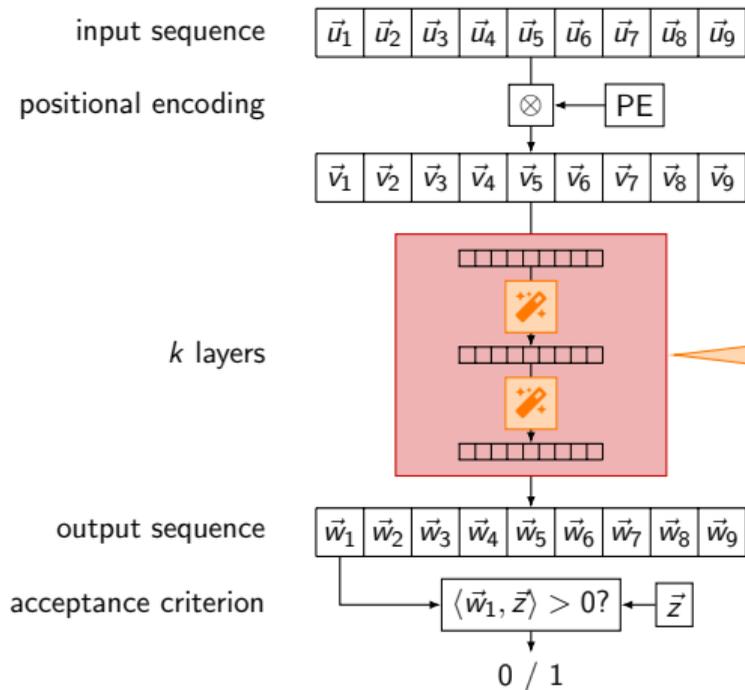
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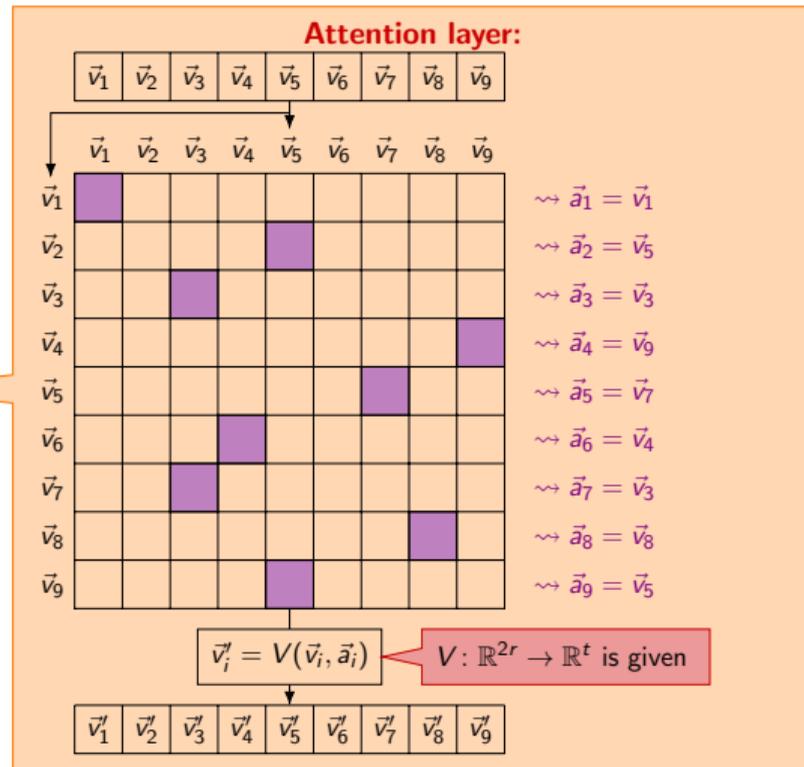
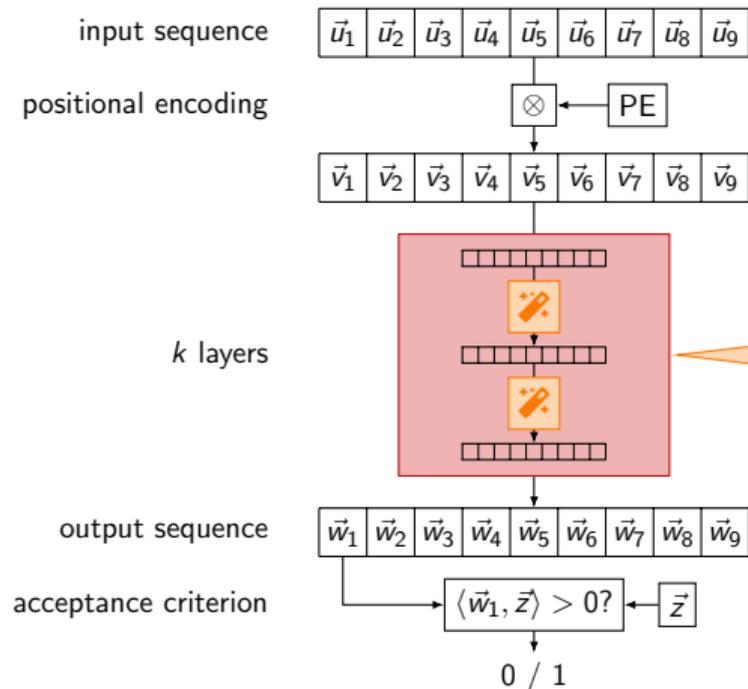
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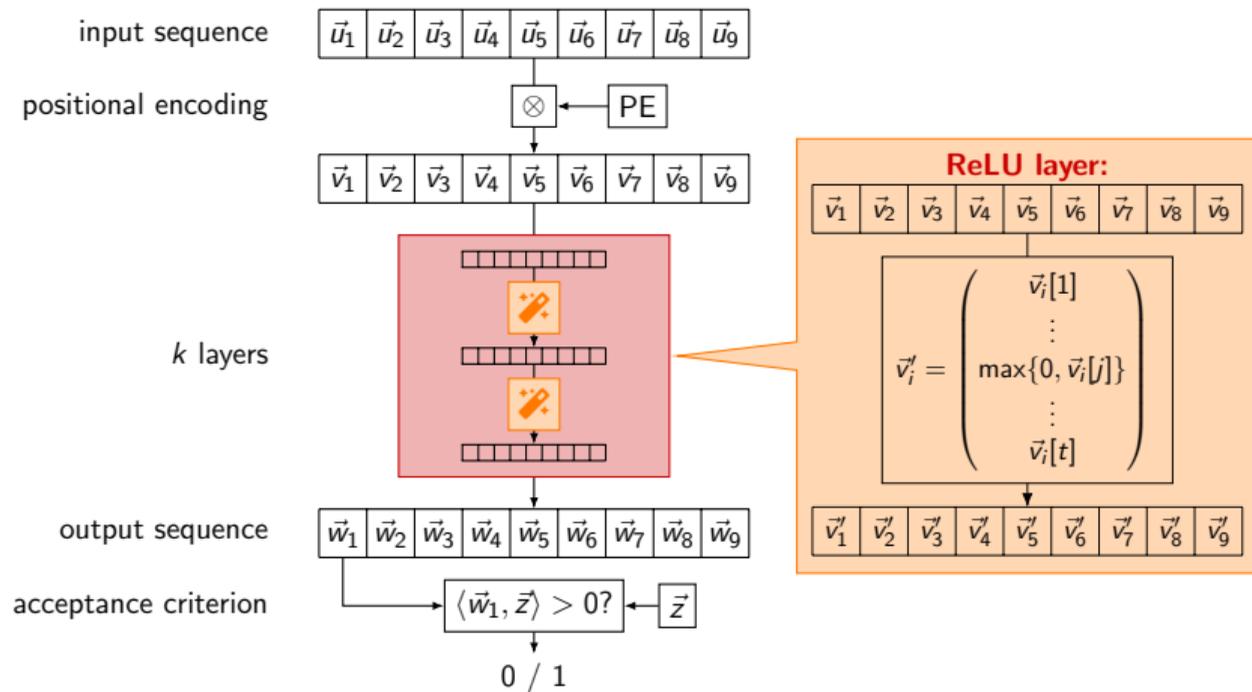
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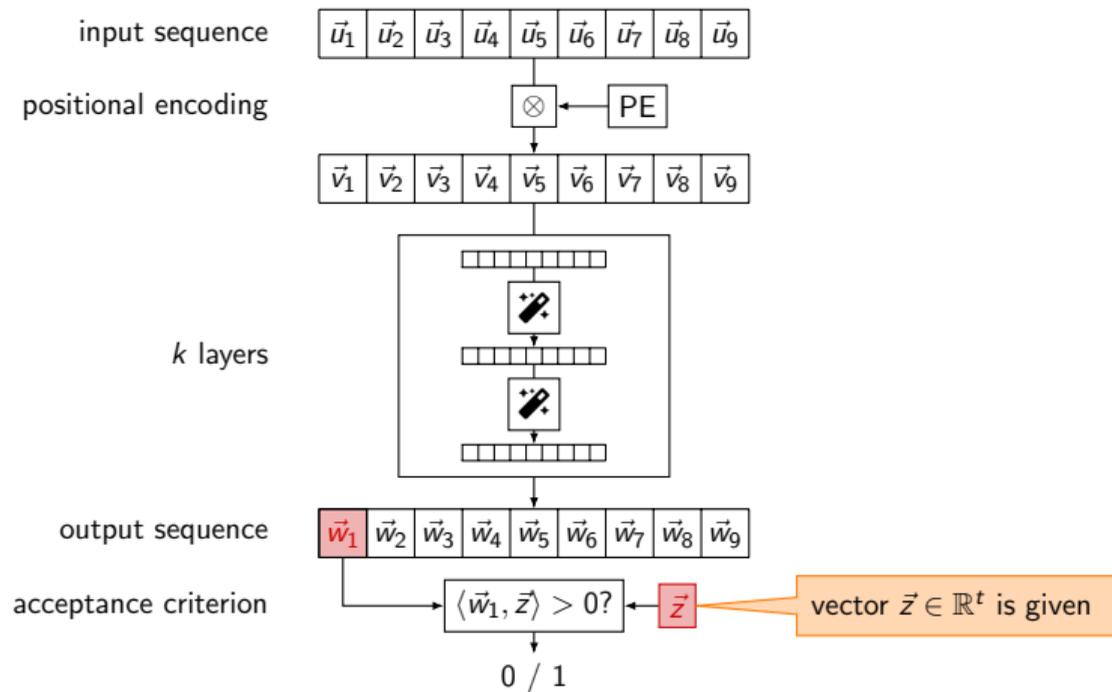
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Word vs. Data Languages

	UHATs with finite alphabet	UHATs with infinite alphabet
Circuit complexity	$\subseteq AC^0$ [Hao et al. 2022] $\neq AC^0$ [Barceló et al. 2023]	$\subseteq TC^0$, $\not\subseteq AC^0$ *
Regularity	$=$ Star-Free * [Angluin et al. 2024]	$\not\subseteq$ Regular *
Logic	$=$ LTL * [Angluin et al. 2024] \supseteq LTL(Mon) [Barceló et al. 2023]	\supseteq LTLTL

Thank you!

Read the paper:



* without positional encoding, but with masking

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Regularity	$= \text{Star-tree}^*$ [Angluin et al. 2024]	$\not\subseteq \text{Regular}^*$
	$\supseteq \text{LTLTL}$	

Languages accepted by a family of circuits of

- constant depth,
- polynomial size, and
- Boolean gates with unbounded fan-in.

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Regularity	= Star-Free * [Angluin et al. 2024]	AC ⁰ + majority gates
Logic	= LTL * [Angluin et al. 2024] $\supseteq LTL(\text{Mon})$ [Barceló et al. 2023]	$\supseteq LTLTL$

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Regular	LTL + monadic predicates Θ [et al. 2024]	$\not\subseteq Regular^*$
Logic	$= L \vee L^*$ [Angluin et al. 2024] $\supseteq LTL(Mon)$ [Barceló et al. 2023]	$\supseteq LTLTL$

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Circuit complexity	$\subseteq AC^0$ [Hao et al. 2022] $\neq AC^0$ [Barceló et al. 2023]	$\subseteq TC^0$, $\not\subseteq AC^0^*$
Regularity	= Sta	Locally Testable LTL = LTL(Mon) + arithmetic constraints
Logic	= LTL* [Angluin et al. 2024] \supseteq LTL(Mon) [Barceló et al. 2023]	\supseteq LTLTL

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