

# Dynamic Trace Estimation

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# Implicit trace estimation

- A basic problem in linear algebra

Given matrix  $A \in \mathbb{R}^{n \times n}$  and access to  $A$  via [matrix vector products](#)

$Av \in \mathbb{R}^n$ , for  $v \in \mathbb{R}^n$  (implicit access), approximate  $\text{tr}(A) = \sum_{i=1}^n A_{ii}$ .

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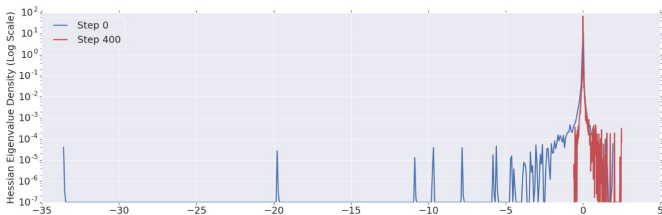


Figure 1: Ghorbani et al. [2019] analyze spectrum of Hessian for Resnet-32.

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- For matrix functions  $A = f(B)$ , we can leverage iterative methods to approximate  $Av = f(B)v$ . (e.g.  $A = B^{-1} / \exp(B) / \log(B)$ ). Typically, runtime is  $O(n^2)$  compared to  $O(n^3)$  for explicitly forming  $A$ .

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- Measure the computational cost in **number of matrix-vector products** required  $Av_1, \dots, Av_\ell$ .

## Hutchinson's estimator (Hutchinson [1990], Girard [1987])

- Approximate  $\text{tr}(A)$  as  $h_\ell(A) = \frac{1}{\ell} \sum_{i=1}^{\ell} g_i^T A g_i$  where entries in  $g_i \in \mathbb{R}^n$  are random i.i.d.  $\pm 1$ .

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- $h_\ell(A)$  approximates  $\text{tr}(A)$  in expectation.

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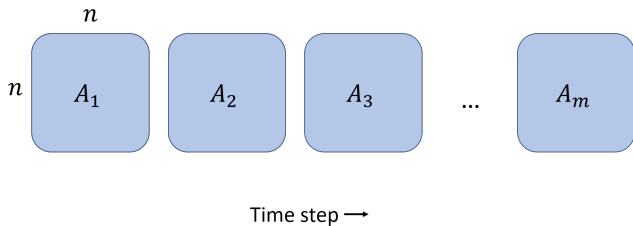
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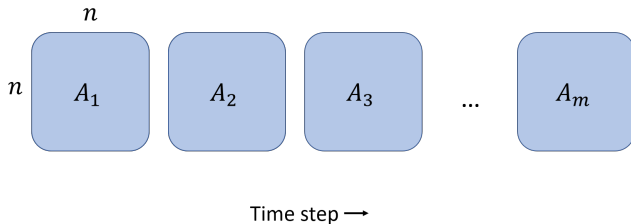
- Variance of  $h_\ell(A) \leq \frac{2}{\ell} \|A\|_F^2$
- For  $\ell = O(\frac{\log(1/\delta)}{\epsilon^2})$ , with high probability,  $|h_\ell(A) - \text{tr}(A)| \leq \epsilon \|A\|_F$

# Dynamic setting



Want good approximations  $t_1, t_2, \dots, t_m$  across all time steps.

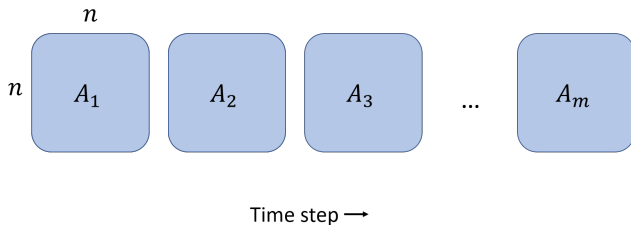
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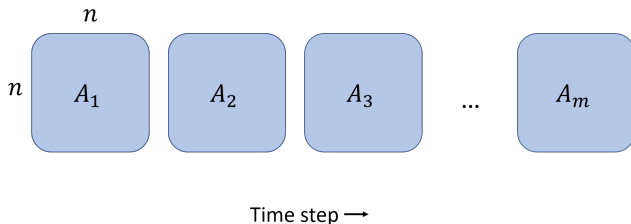


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**Our result:** Yes and can obtain quadratic improvements under certain assumptions!



## Problem (Dynamic trace estimation)

Let  $A_1, \dots, A_m$  be  $n \times n$  symmetric matrices satisfying:

1.  $\|A_i\|_F \leq 1$ , for all  $i \in [1, m]$ .
2.  $\|A_{i+1} - A_i\|_F \leq \alpha$ , for all  $i \in [1, m - 1]$ .

Given implicit matrix-vector multiplication access to each  $A_i$  in sequence, the goal is to compute trace approximations  $t_1, \dots, t_m$  for  $\text{tr}(A_1), \dots, \text{tr}(A_m)$  such that, for each  $i \in 1, \dots, m$ ,

$$\mathbb{P}[|t_i - \text{tr}(A_i)| \geq \epsilon] \leq \delta.$$

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DeltaShift:  $t_{i+1} = (1 - \gamma)t_i + h_\ell(A_{i+1} - (1 - \gamma)A_i)$

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- Multiplying by  $(1 - \gamma)$  reduces the variance of the leading term.

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For any  $\epsilon, \delta, \alpha \in (0, 1)$ , the DeltaShift algorithm solves Dynamic Trace Estimation problem with

$$O\left(m \cdot \frac{\alpha \log(1/\delta)}{\epsilon^2} + \frac{\log(1/\delta)}{\epsilon^2}\right)$$

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For  $\alpha \approx \epsilon$ , DeltaShift requires  $O\left(\frac{\log(1/\delta)}{\epsilon}\right)$  total matrix-vector products.

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- Note: We can reuse the same matrix-vector products used by trace estimation.

For a PSD matrix, recent algorithm by Meyer et al. [2021] obtains the  $(\epsilon, \delta)$  bounds with  $\frac{\log(1/\delta)}{\epsilon}$  matrix-vector products.

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For stronger assumptions (in form of nuclear norm) on sequence of matrices:

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For  $\|A_i\|_* \leq 1$  and  $\|A_{i+1} - A_i\|_* \leq \alpha$  for all  $i$ , DeltaShift++ solves dynamic trace estimation problem with

$$O\left(m \cdot \frac{\sqrt{\alpha/\delta}}{\epsilon} + \frac{\sqrt{1/\delta}}{\epsilon}\right)$$

total matrix-vector products with  $A_1, A_2, \dots, A_m$ .



We can estimate near-optimal  $\gamma$  for DeltaShift++ as well!

Let  $K_A = \|A - A_k\|_F^2$

$$\gamma_j^* = \min_{\gamma} \left[ \frac{\gamma^2 8K_{A_j}}{\ell} + (1 - \gamma)^2 \left( v_{j-1} + \frac{8K_{\Delta_j}}{\ell} \right) \right]$$

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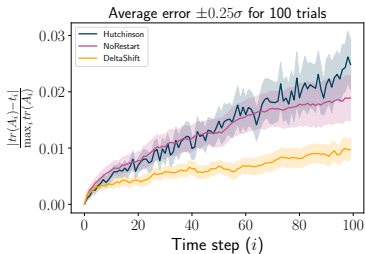
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Similar to DeltaShift, we can reuse matrix-vector products from trace estimation!

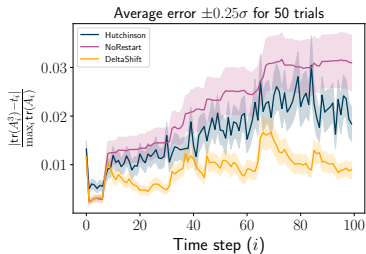
For the dynamic trace problem, we compare using the **same number of total matrix-products** for

- Hutchinson's estimator at each time step
- Estimate  $\text{tr}(\Delta_i)$  at each time step and add to  $\text{tr}(A_i)$  (NoRestart)
- DeltaShift

# Empirical results



(a) Synthetic data with total matrix-vector products =  $8 * 10^3$



(b) Graph data with total matrix-vector products =  $10^4$

- For estimating spectral density, trace of polynomials of the Hessian is used.

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The three term recurrence relation for Chebyshev polynomials is:

$$T_0(H) = I, \quad T_1(H) = H, \quad T_{n+1}(H) = 2HT_n(H) - T_{n-1}(H).$$

**Table 1:** Average error for trace of polynomials of Hessian with learning rate 0.001 and total matrix-vector products = 2000

	HUTCHINSON	NORESTART	DELTA SHIFT
$T_1(H)$	2.5E-02	3.7E-02	<b>1.7E-02</b>
$T_2(H)$	1.2E-06	1.7E-06	<b>8.0E-07</b>
$T_3(H)$	4.0E-02	4.1E-02	<b>3.1E-02</b>
$T_4(H)$	1.5E-06	1.7E-06	<b>1.0E-06</b>
$T_5(H)$	2.1E-02	4.3E-02	<b>1.9E-02</b>



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- Can we do better when  $\Delta$  matrices have additional structure?  
Partial progress in form of DeltaShift++.

Thank you!